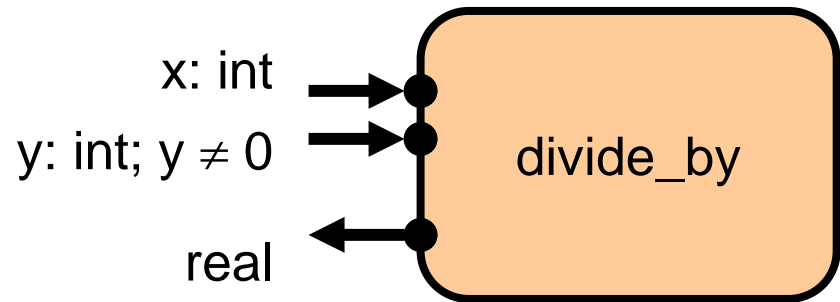# Interface-based Design 3

Tom Henzinger
EPFL and UC Berkeley
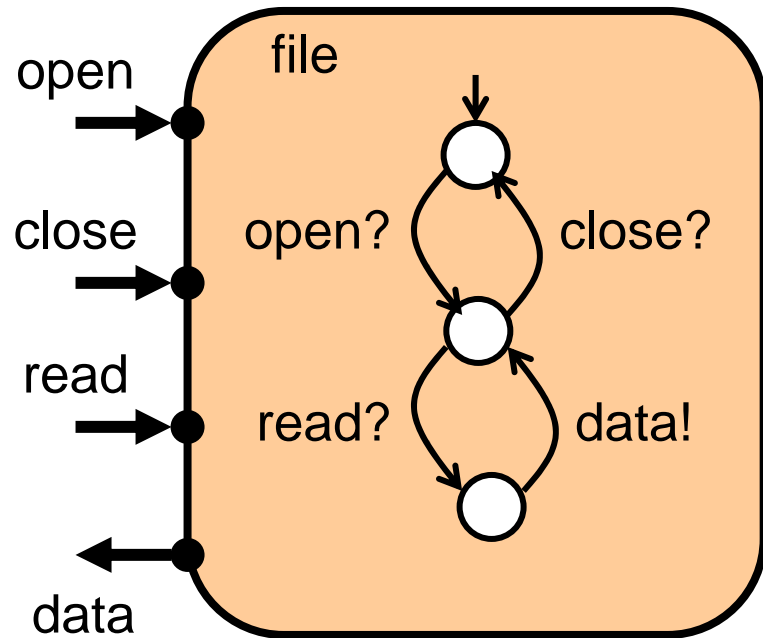
An Assertional Interface

This interface constrains the client's data.
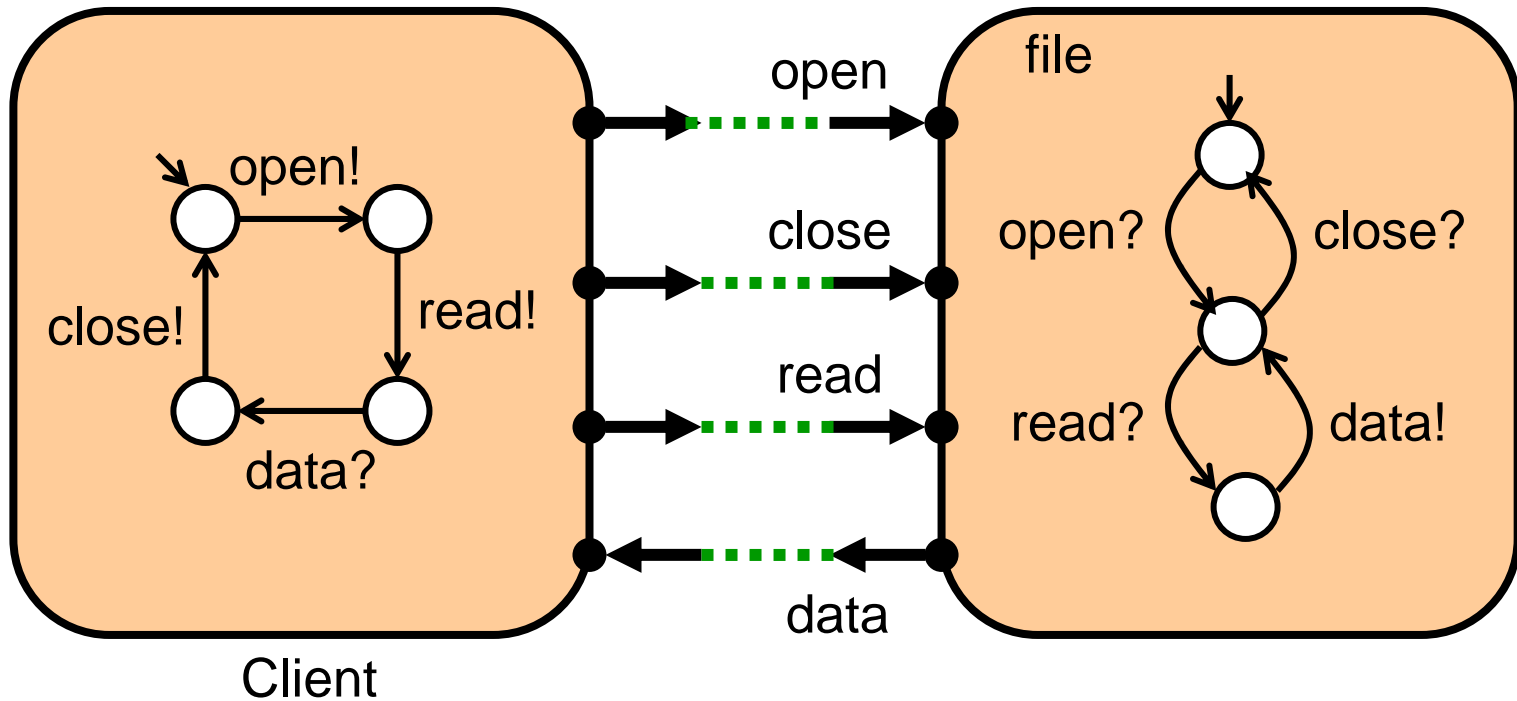

x: int

y: int; $y \neq 0$

real

divide_by

An Automaton Interface
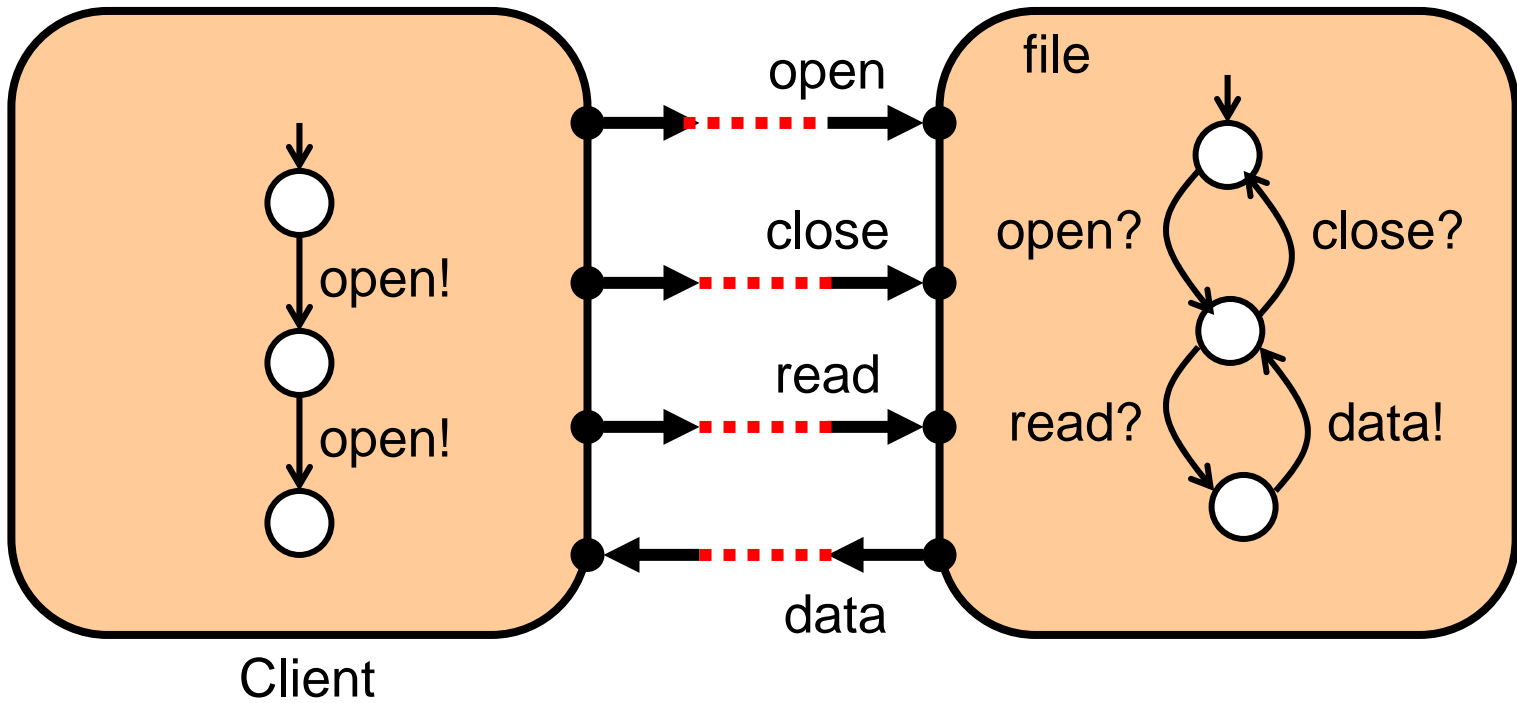
This interface
constrains the
client's control.

Two Compatible Automaton Interfaces

Two Incompatible Automaton Interfaces

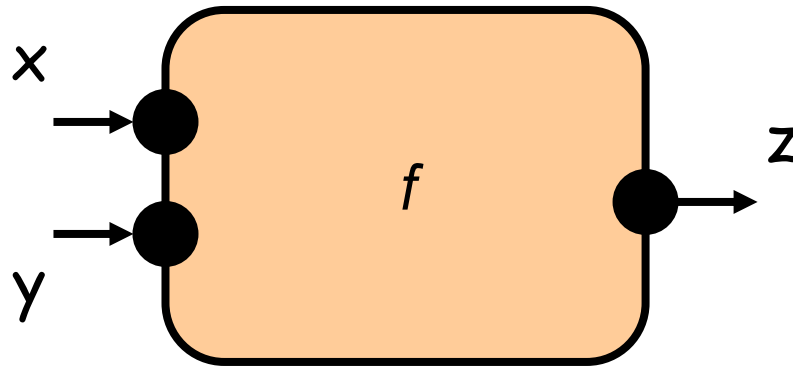*Today's Lecture:*

How do we check the compatibility of automaton interfaces ?

What is the composition of two compatible automaton interfaces ?

Free Inputs:
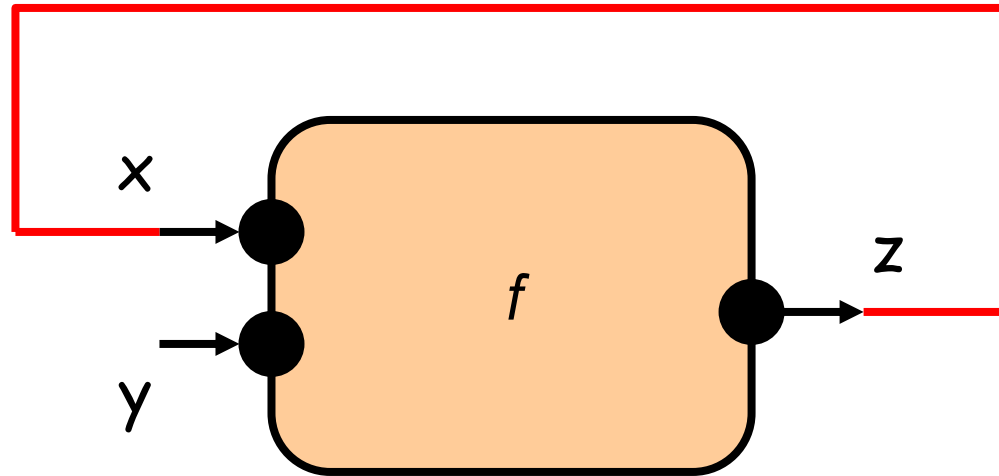*Interface Composition propagates Environment Constraints.*

$$x=0 \implies y=0 \qquad\qquad true$$

Free Inputs:
Interface Composition propagates Environment Constraints.

$x=0 \Rightarrow y=0$ true

Free Inputs:
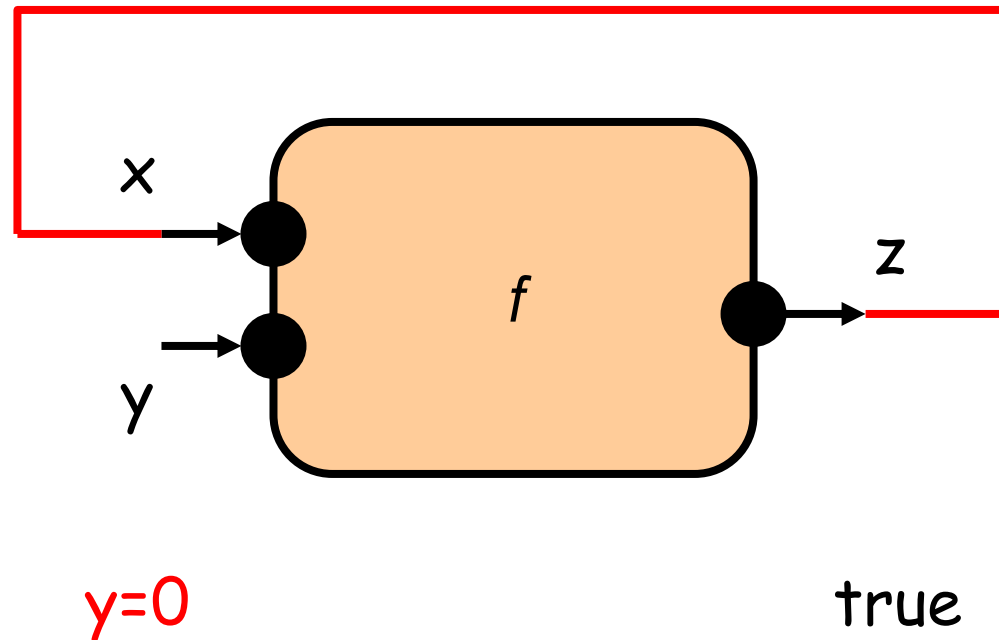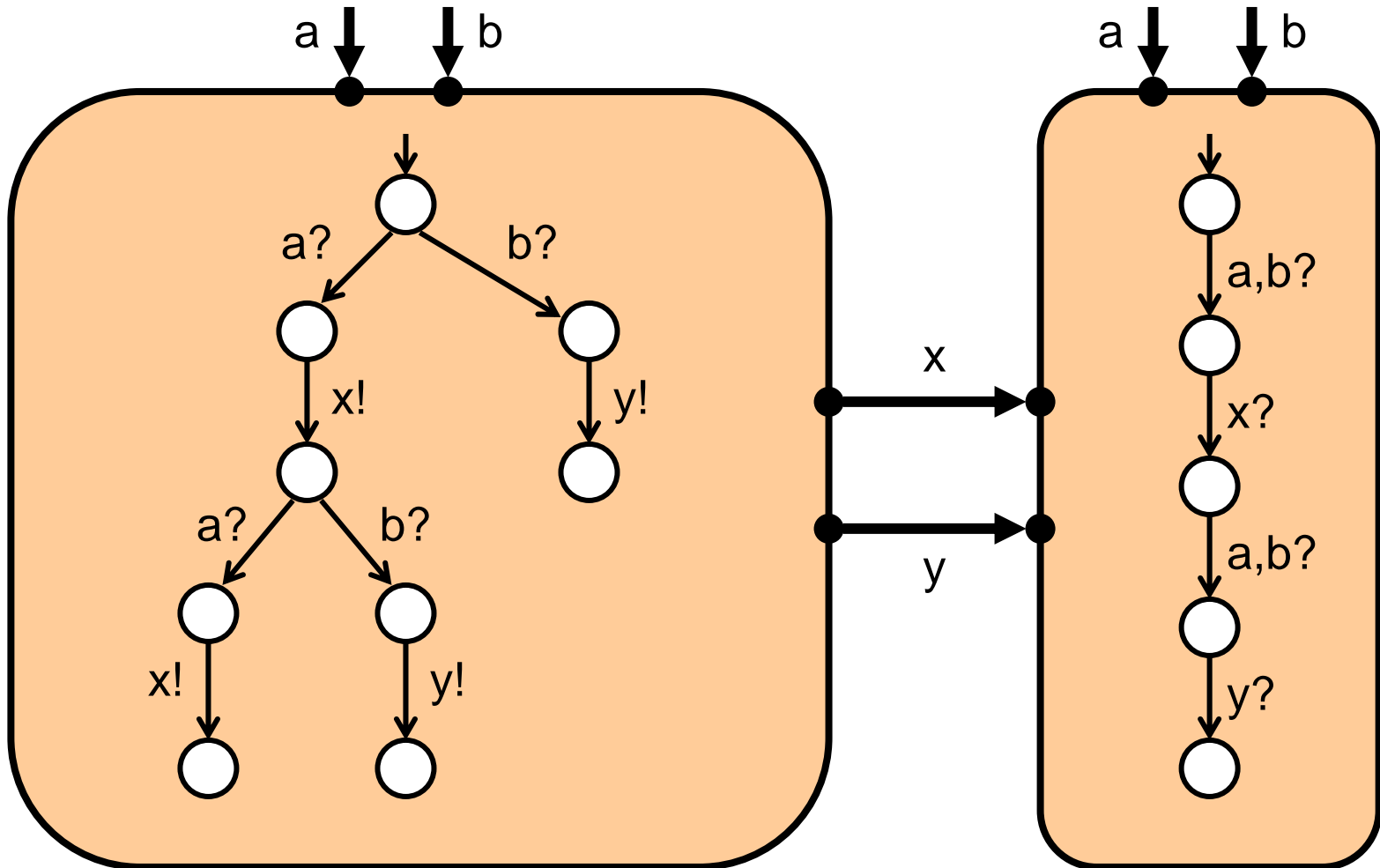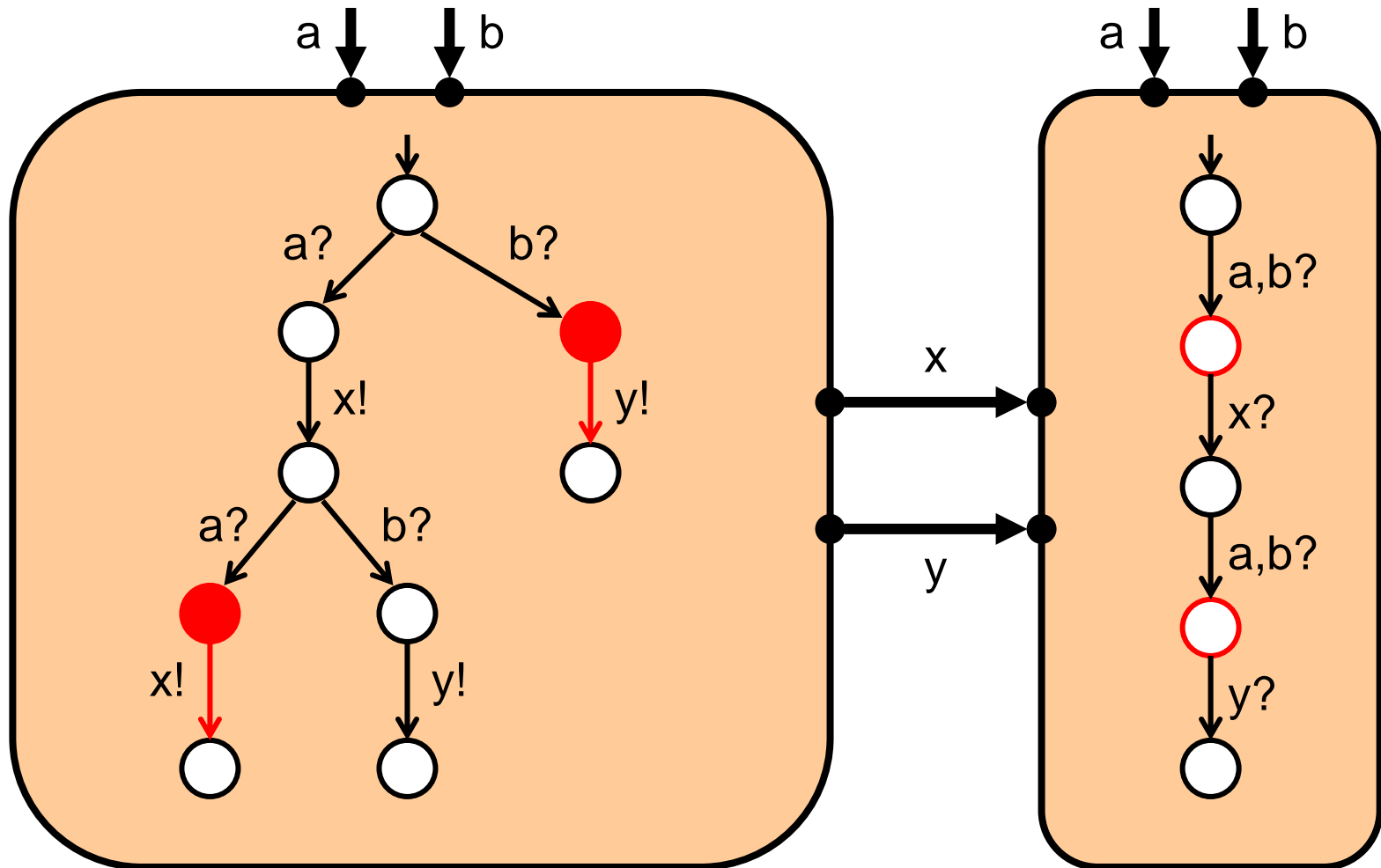## *Interface Composition propagates Environment Constraints.*
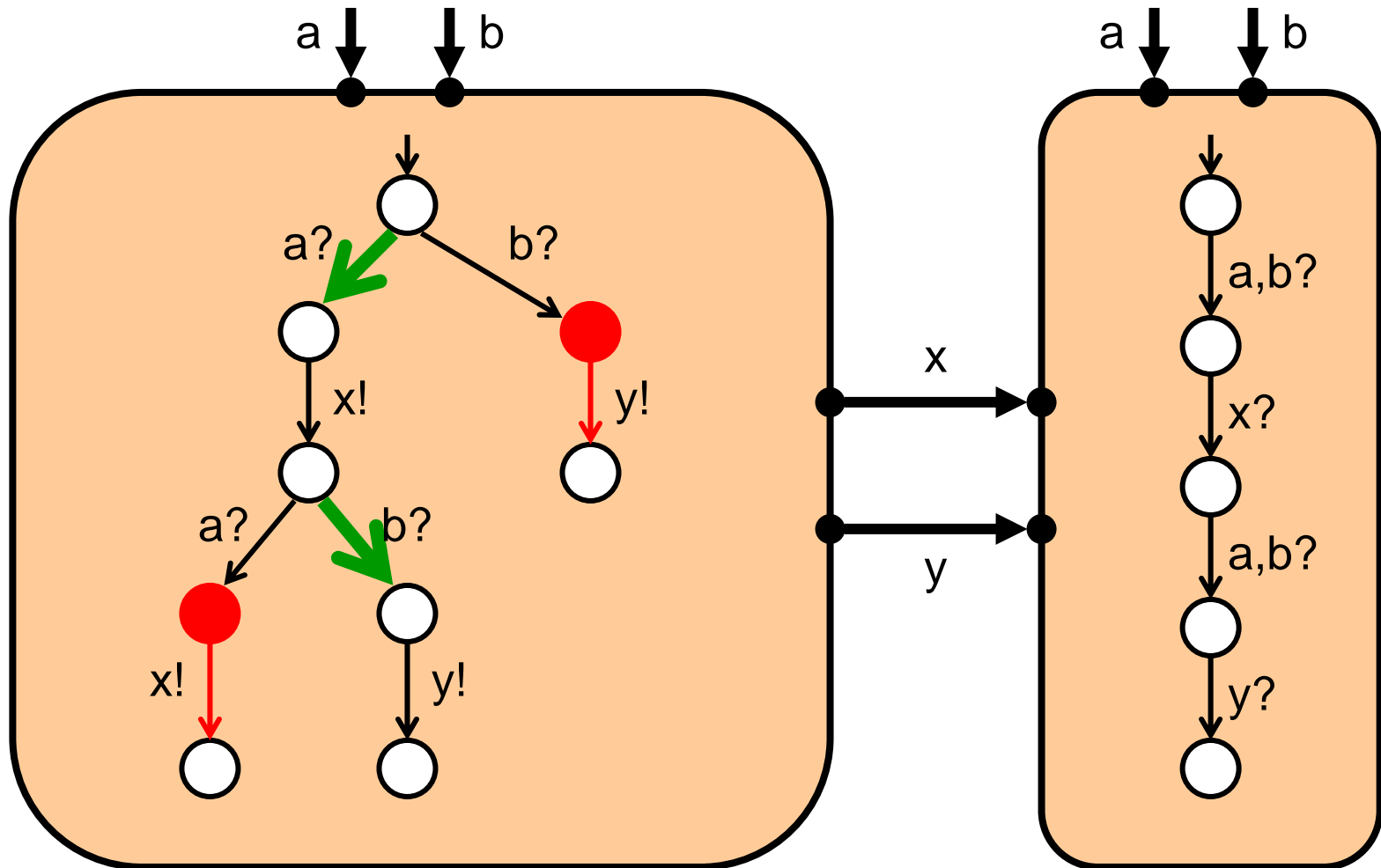

y=0

true

The environment is helpful !
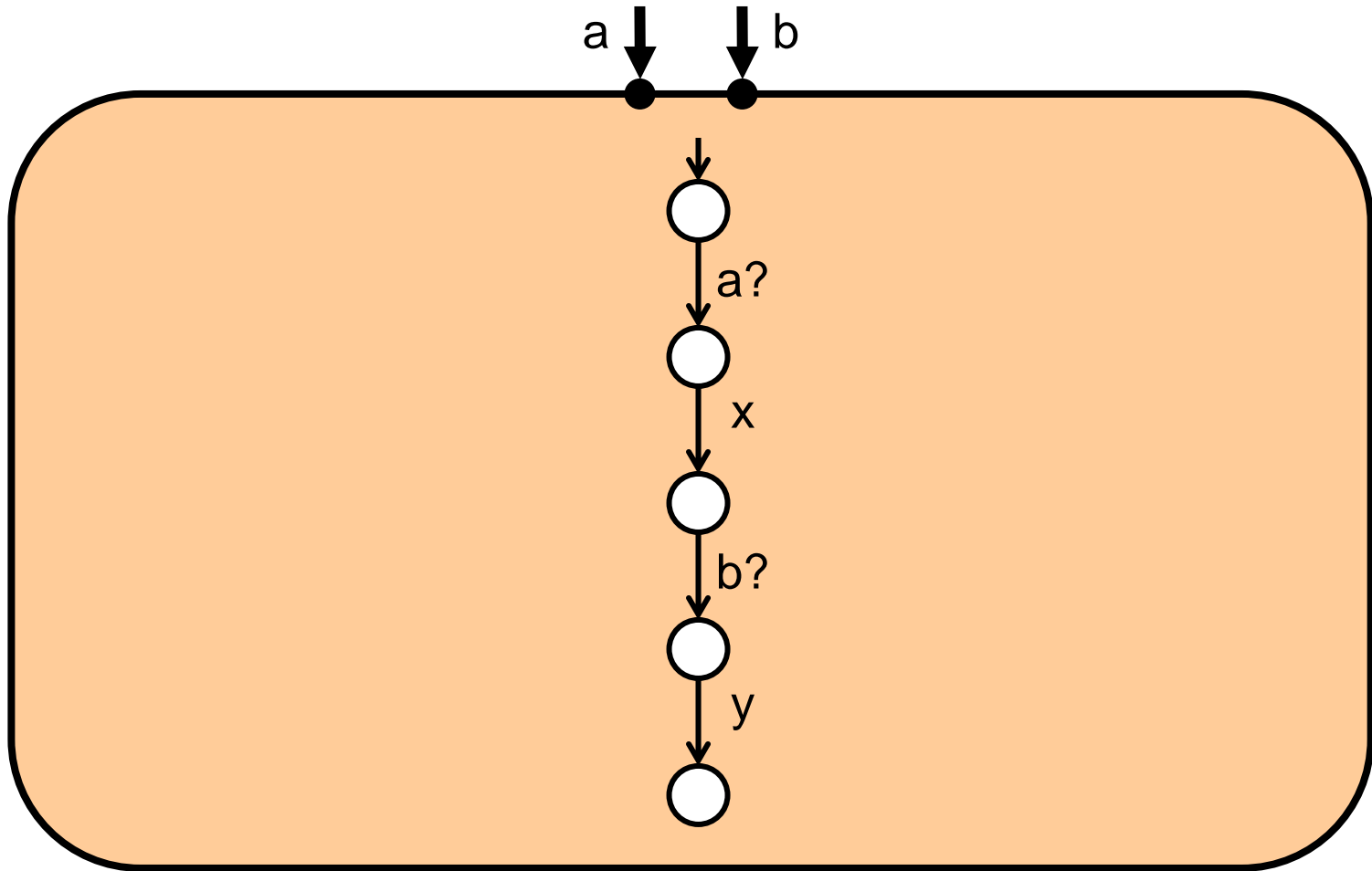
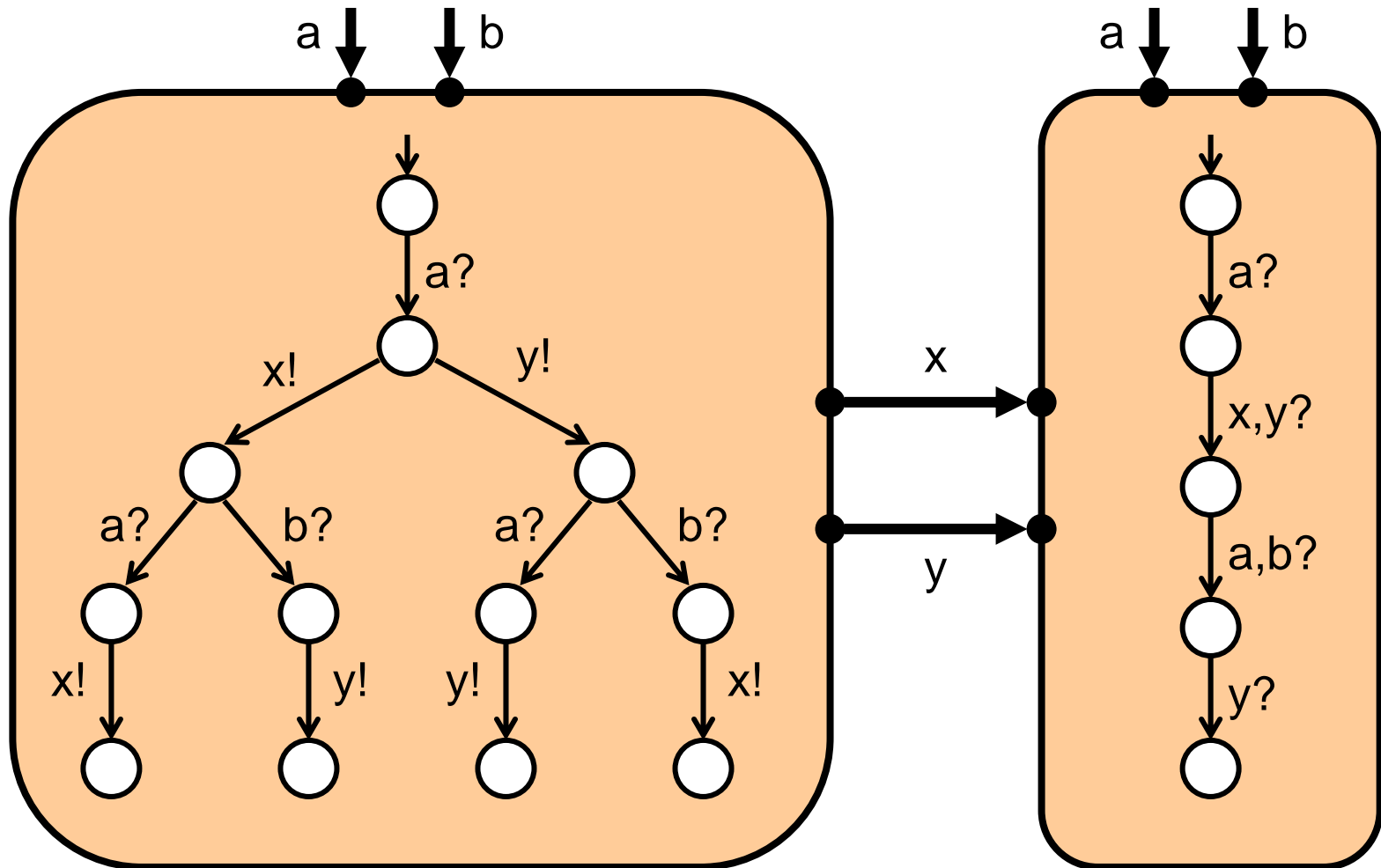# Automaton Interface Composition

Automaton Interface Composition

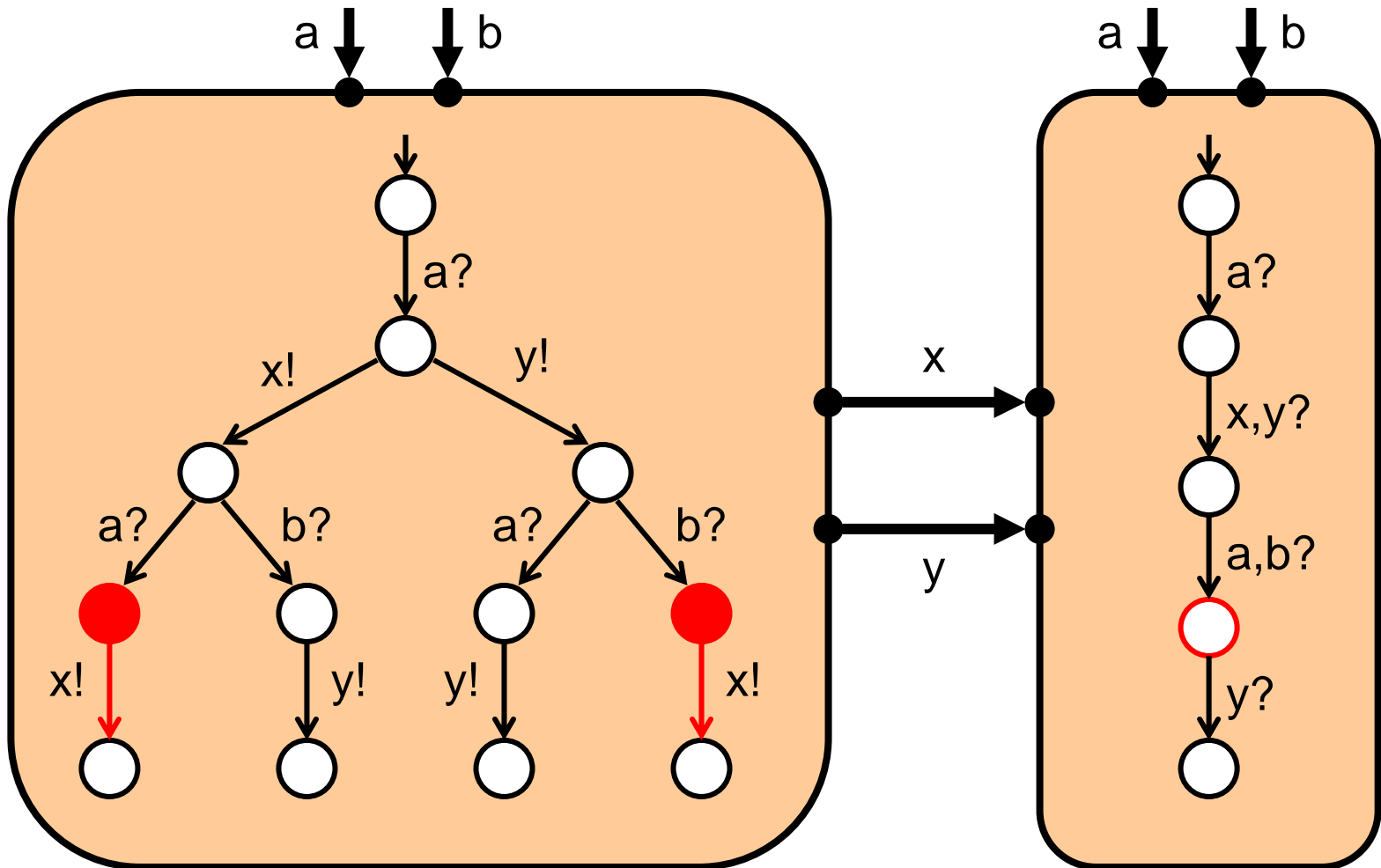# Automaton Interface Composition

The Composite Interface

Automaton Interface Composition

Automaton Interface Composition

Automaton Interface Composition

The Composite Interface

Lesson 3:

Stateful Interfaces are Games!

-Player Input vs. player Internal.

-The composite interface is the product restricted to those states from which player Input has a strategy to avoid incompatibilities.

Graph Games

AND-OR Graph:

OR Nodes
AND Nodes

Graph Games

AND-OR Graph:

OR Player
AND Player

Graph Games

AND-OR Graph:

OR Player
AND Player

Graph Games

AND-OR Graph:

OR Player
AND Player

Graph Games

AND-OR Graph:

OR Player
AND Player

Safety Games

AND-OR Graph:

OR Player
AND Player

From which nodes does the OR Player have a strategy to avoid ERROR Nodes ?
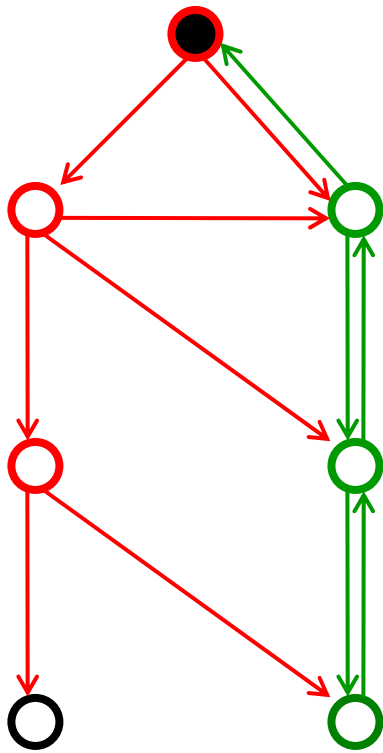
Safety Games

AND-OR Graph:

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Safety Games

AND-OR Graph:

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From which nodes does the
OR Player have a strategy
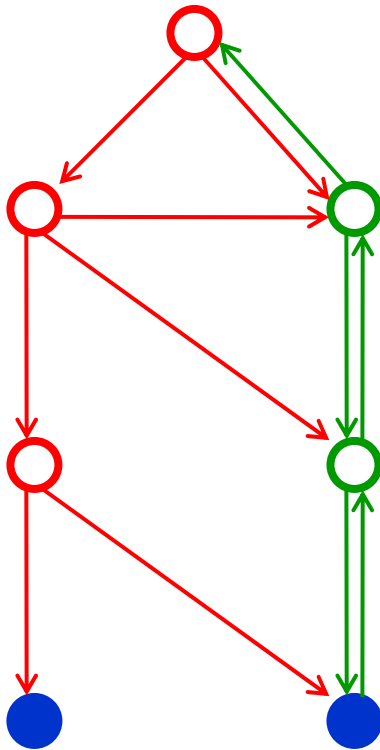to avoid ERROR Nodes ?

Safety Games

AND-OR Graph:

OR Player
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From which nodes does the OR Player have a strategy to avoid ERROR Nodes ?
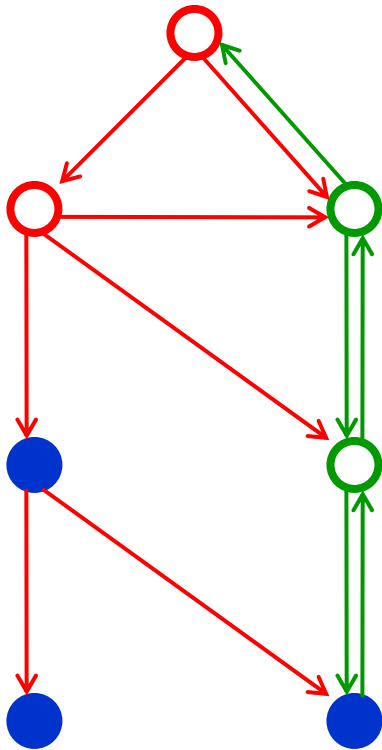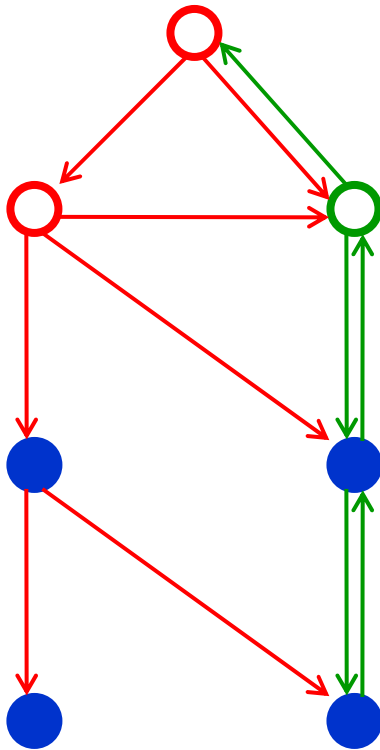
Complexity?

Safety Games

AND-OR Graph:

OR Player
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From which nodes does the
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Safety Games

AND-OR Graph:

OR Player
AND Player

From which nodes does the
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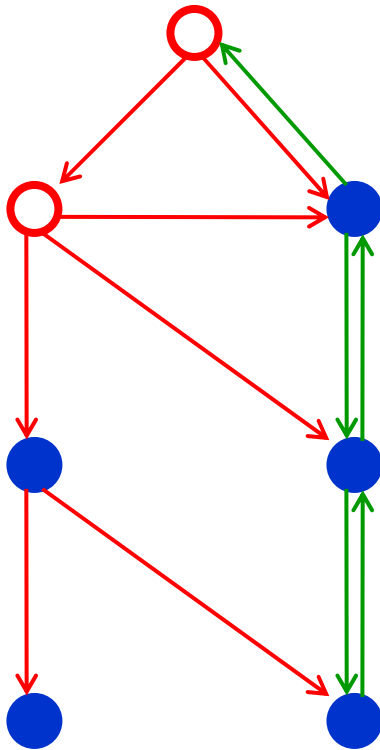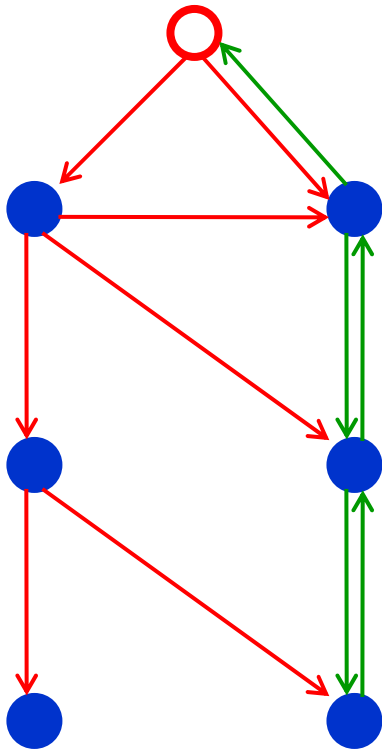to avoid ERROR Nodes ?

a b

Safety Games

AND-OR Graph:

OR Player
AND Player

From which nodes does the OR Player have a strategy to avoid ERROR Nodes ?

b

Safety Games

AND-OR Graph:

OR Player
AND Player

From which nodes does the
OR Player have a strategy
to avoid ERROR Nodes ?

b c

Safety Games

AND-OR Graph:

OR Player
AND Player

From which nodes does the
OR Player have a strategy
to avoid ERROR Nodes ?

b c

Safety Games

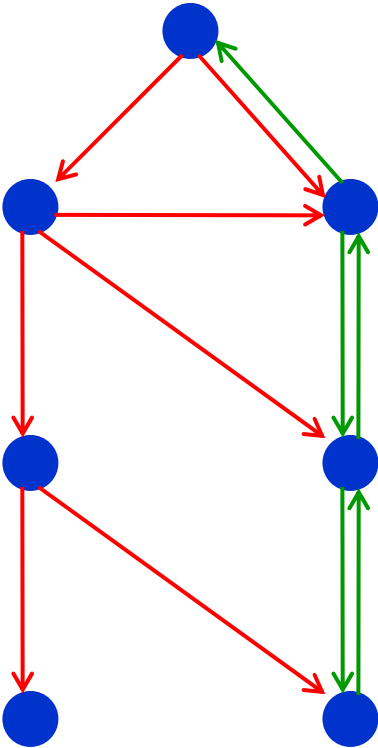AND-OR Graph:

OR Player
AND Player

From which nodes does the OR Player have a strategy to avoid ERROR Nodes ?
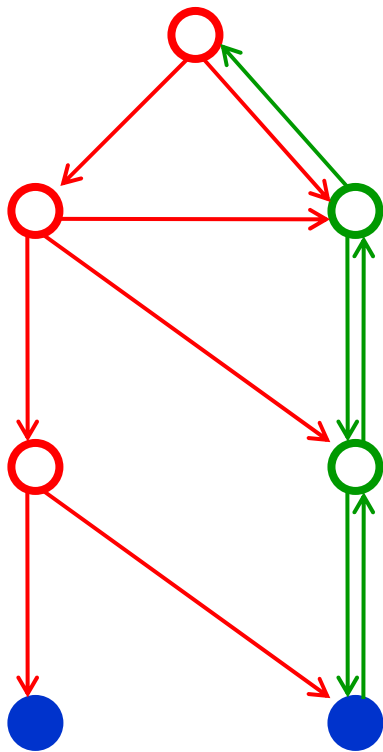
c

Safety Games

AND-OR Graph:

OR Player
AND Player

From which nodes does the OR Player have a strategy to avoid ERROR Nodes ?

c d

Safety Games

AND-OR Graph:

OR Player
AND Player

From which nodes does the
OR Player have a strategy
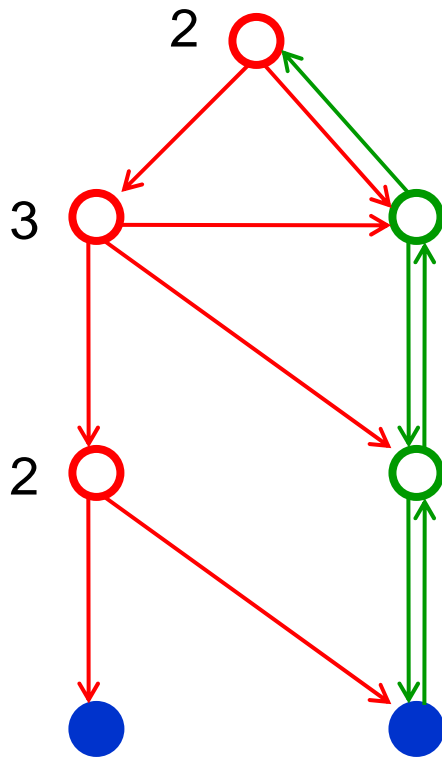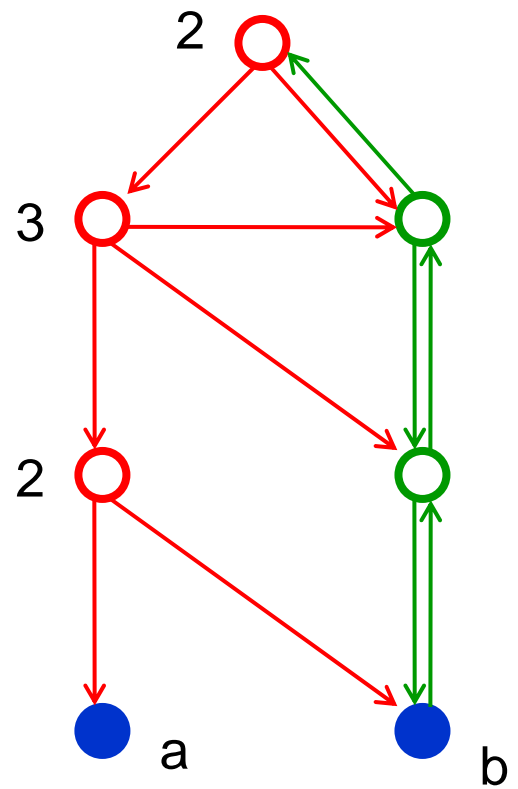to avoid ERROR Nodes ?

d

Safety Games

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AND Player

From which nodes does the OR Player have a strategy to avoid ERROR Nodes ?

d

Safety Games

AND-OR Graph:

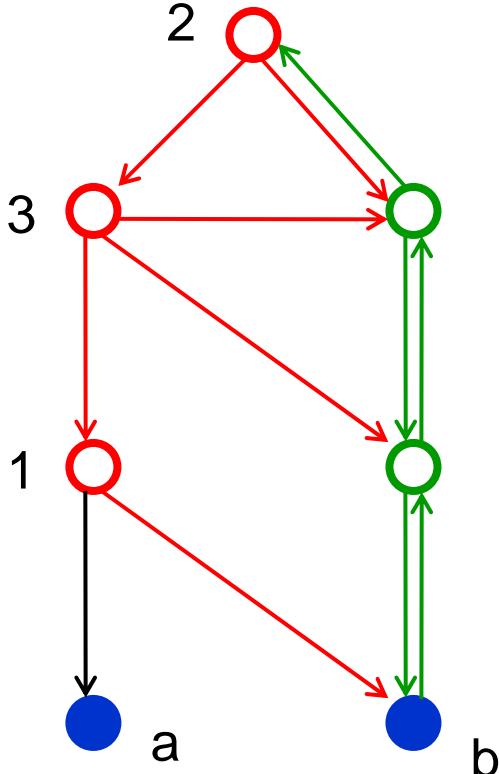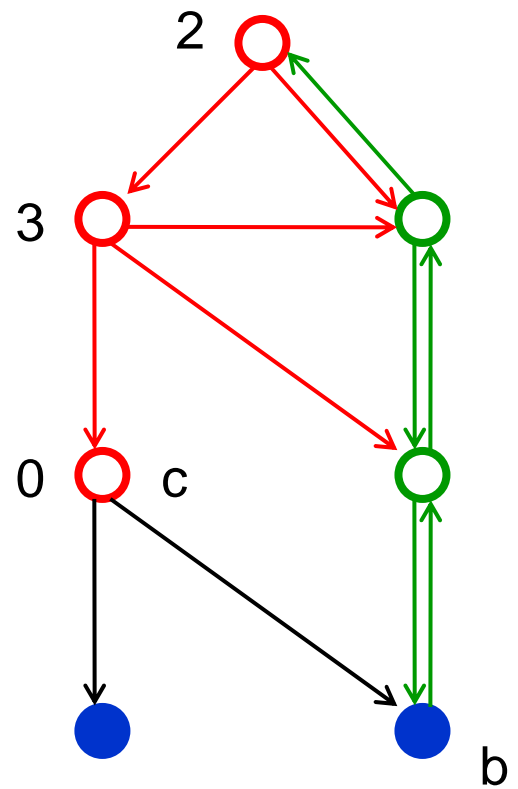OR Player
AND Player

From which nodes does the
OR Player have a strategy
to avoid ERROR Nodes ?

d

Safety Games

AND-OR Graph:

OR Player
AND Player

From which nodes does the
OR Player have a strategy
to avoid ERROR Nodes ?

Safety Games

2

1   e

AND-OR Graph:
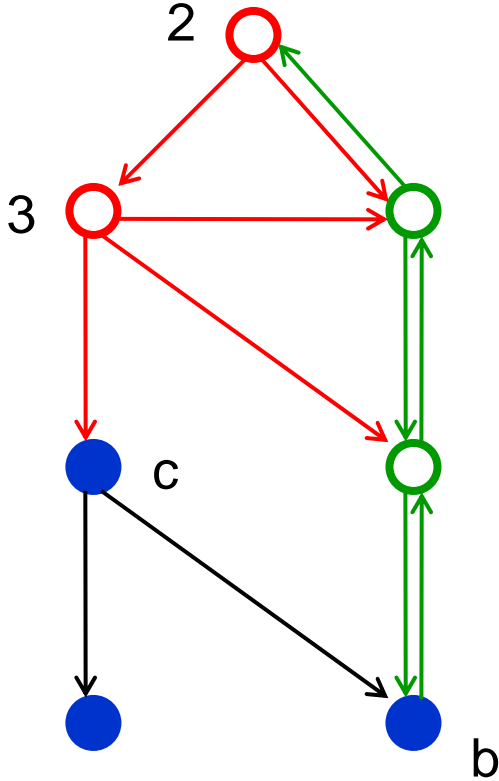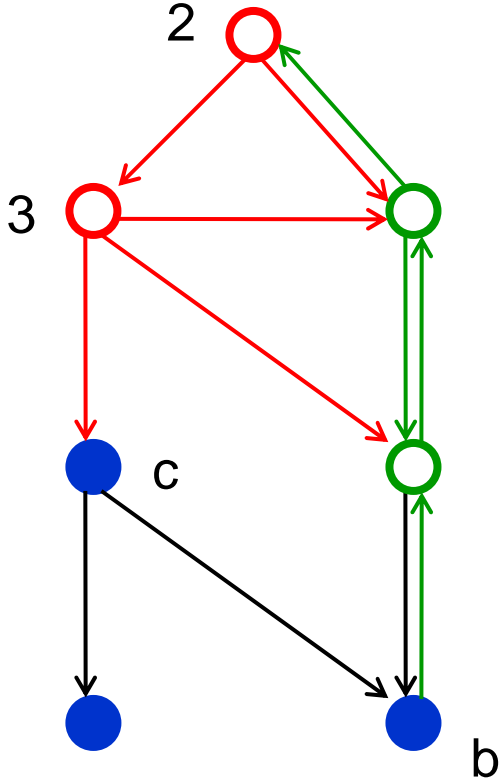
OR Player
AND Player

From which nodes does the
OR Player have a strategy
to avoid ERROR Nodes ?

e

Safety Games

AND-OR Graph:

OR Player
AND Player

From which nodes does the
OR Player have a strategy
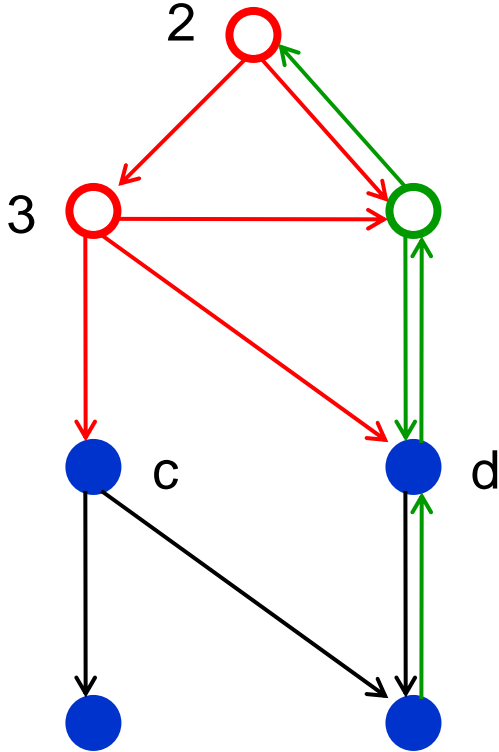to avoid ERROR Nodes ?

e

Safety Games

AND-OR Graph:

OR Player
AND Player

From which nodes does the
OR Player have a strategy
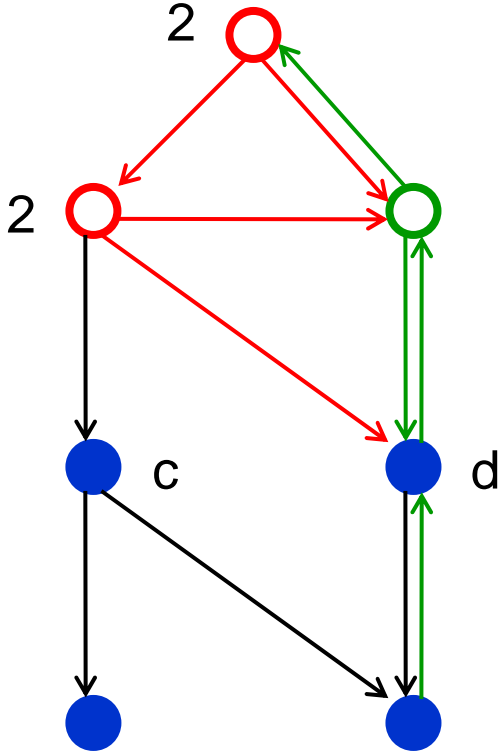to avoid ERROR Nodes ?

e

Safety Games

AND-OR Graph:

OR Player
AND Player

From which nodes does the
OR Player have a strategy
to avoid ERROR Nodes ?

e f

Safety Games

AND-OR Graph:

OR Player
AND Player

From which nodes does the OR Player have a strategy to avoid ERROR Nodes ?
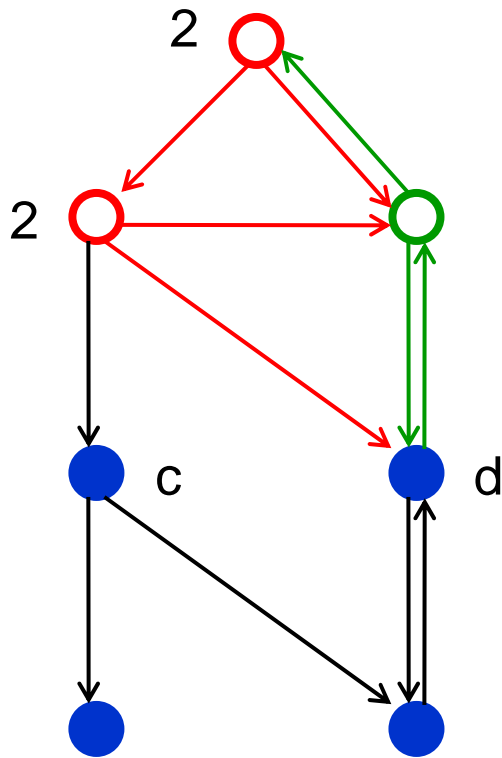
f

Safety Games

AND-OR Graph:

OR Player
AND Player

From which nodes does the
OR Player have a strategy
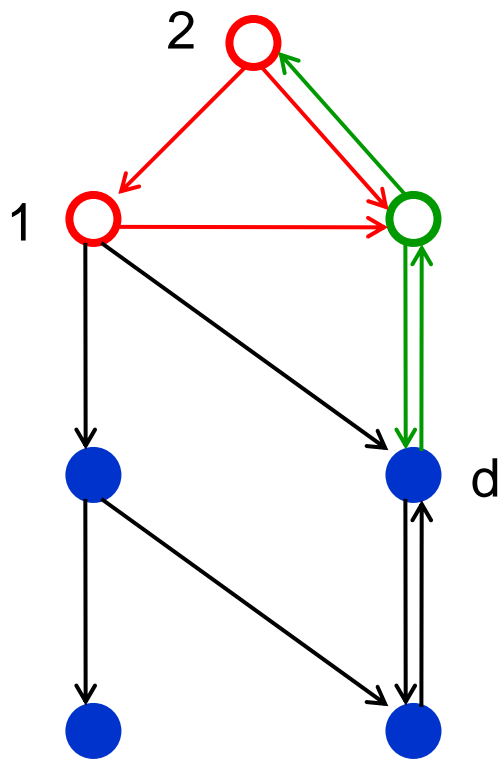to avoid ERROR Nodes ?

Safety Games

AND-OR Graph:

OR Player
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From which nodes does the
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to avoid ERROR Nodes ?

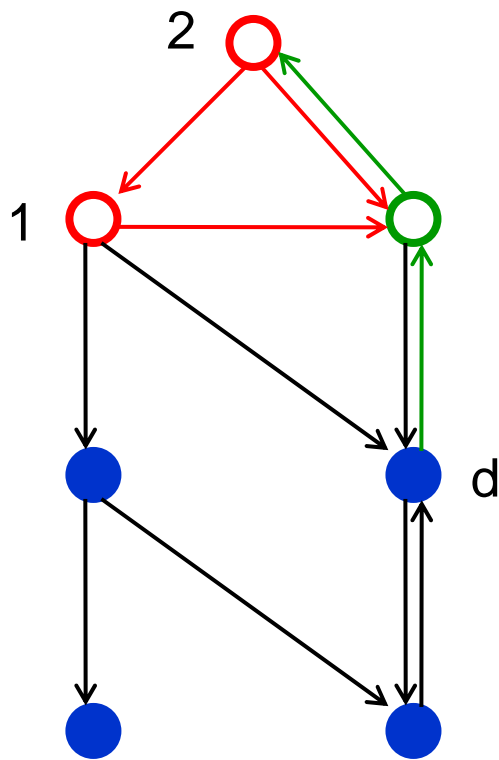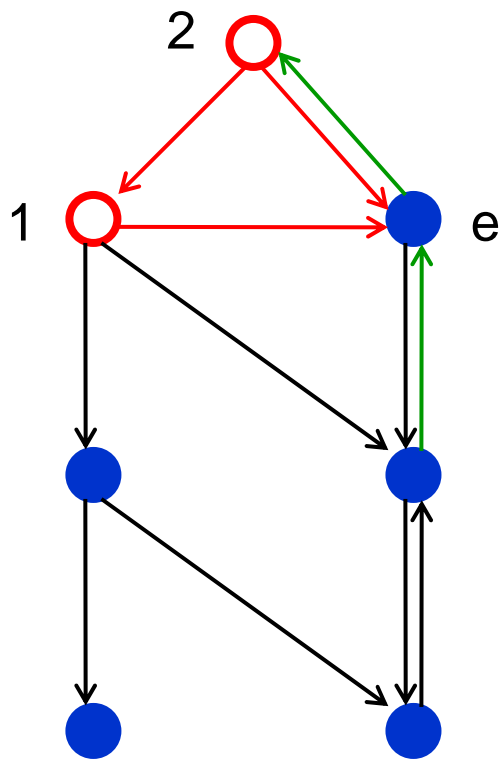Linear Time (P-Complete)

Safety Games

AND-OR Graph:

OR Player
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From which nodes does the
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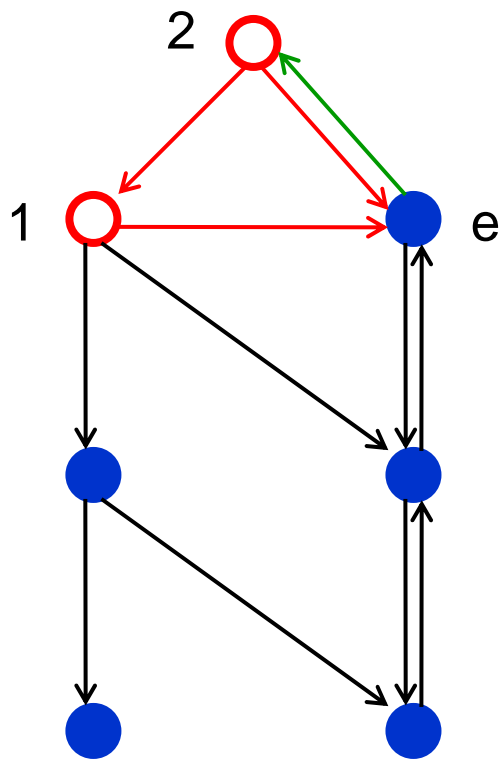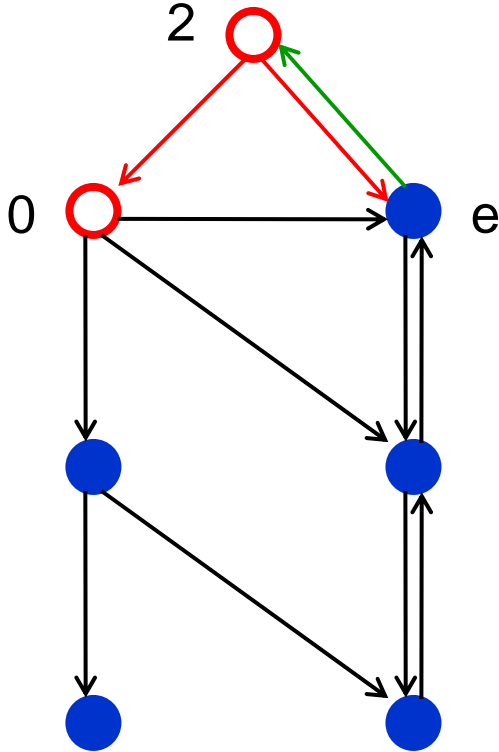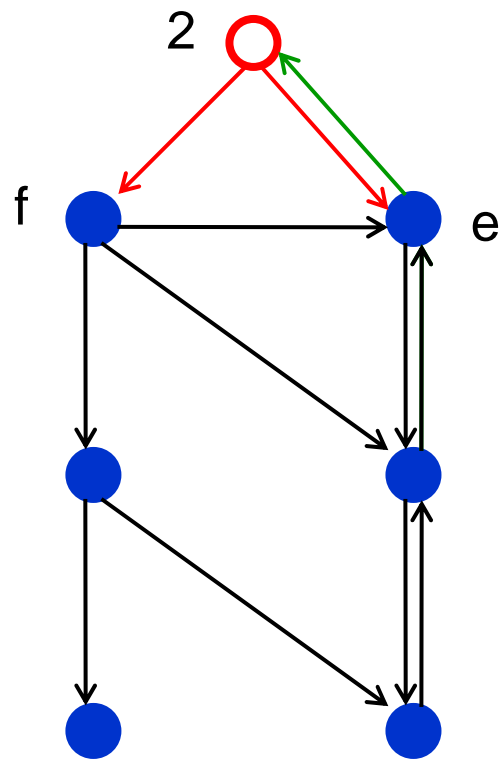to avoid ERROR Nodes ?

Safety Games

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Safety Games

AND-OR Graph:
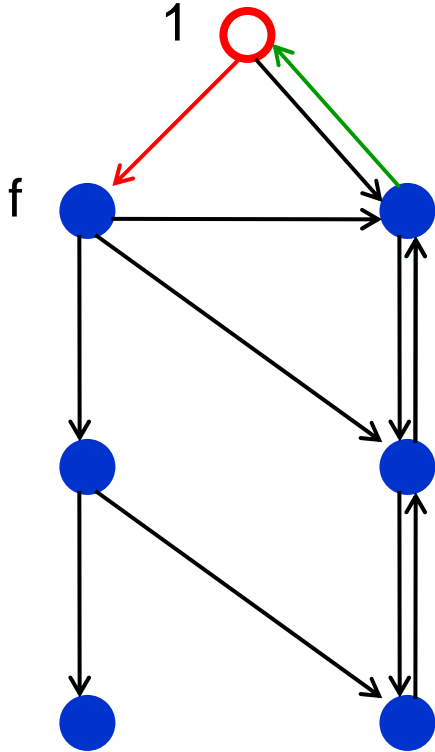
OR Player
AND Player

From which nodes does the
OR Player have a strategy
to avoid ERROR Nodes ?

Most general memoryless
pure strategy exists.

From Interfaces to Games

From Interfaces to Games

ERROR states of product automaton.

From Interfaces to Games

OR Player   ...   External choices made by the environment
AND Player   ...   Internal choices made by the interface product

From Interfaces to Games

Does the Environment have a strategy to avoid the ERROR states?

From Interfaces to Games

Yes

Does the Environment have a strategy to avoid the ERROR states?

From Interfaces to Games

The most general environment strategy.

The Composite Interface

So far, we have used a simple lock-step model of concurrency.

An interface formalism can be built around any model of concurrency.

So far, we have used a simple lock-step model of concurrency.

An interface formalism can be built around any model of concurrency.

Example:

I/O Automata [Lynch]

⬇

Interface Automata [deAlfaro,H]

So far, we have used a simple lock-step model of concurrency.

An interface formalism can be built around any model of concurrency.

Example:

I/O Automata [Lynch]

-Total composition
-A process model

same syntax ⬇ *different semantics*

Interface Automata [deAlfaro,H]

-Partial composition
-Compatibility check
-An interface model

An Interface Automaton

$$F = (Q_F,\ Q^0_F,\ A^I_F,\ A^O_F,\ A^H_F,\ T_F)$$

Q_F ... set of *states*

$Q^0_F \subseteq Q_F$ ... set of *initial states*

A^I_F ... *input actions*

A^O_F ... *output actions*

A^H_F ... *internal* (hidden) *actions*

mutually disjoint
$A_F = A^I_F \cup A^O_F \cup A^H_F$

$T_F \subseteq Q_F \times A_F \times Q_F$ ... set of *transitions*

The Product of Interface Automata

F and G are *composable* if

$$A^H_F \text{ Å } A_G = ; \qquad A^H_G \text{ Å } A_F = ; \qquad A^I_F \text{ Å } A^I_G = ; \qquad A^O_F \text{ Å } A^O_G = ;$$

If F and G are composable, then

$$Q_{F \pounds G} = Q_F \pounds Q_G$$

$$Q^0_{F \pounds G} = Q^0_F \pounds Q^0_G$$

$$A^I_{F \pounds G} = (A^I_F [A^I_G) \setminus shared(F,G)$$

$$A^O_{F \pounds G} = (A^O_F [A^O_G) \setminus shared(F,G)$$

$$A^H_{F \pounds G} = A^H_F [A^H_G [shared(F,G)$$

$$shared(F,G) = A_F \text{ Å } A_G$$

The Product of Interface Automata, continued

$T_{F \pounds G}$ =

$\{ ((f,g), a, (f',g)) : (f,a,f') \, 2 \, T_F \, \text{Æ} \, a \notin shared(F,G) \} \, [$

$\{ ((f,g), b, (f,g')) : (g,b,g') \, 2 \, T_G \, \text{Æ} \, b \notin shared(F,G) \} \, [$

$\{ ((f,g), c, (f',g')) : (f,c,f') \, 2 \, T_F \, \text{Æ} \, (g,c,g') \, 2 \, T_G \, \text{Æ} \, c \, 2 \, shared(F,G) \}$

The product automaton.

The Error States of the Product Automaton

Error(F,G) = { (f,g) :   (9 a 2 shared(F,G))

$$\text{(a 2 } A^O_F(f) \text{ Æ } a \notin A^I_G(g)) \text{ Ç}$$

$$\text{(a 2 } A^O_G(g) \text{ Æ } a \notin A^I_F(f)) \text{ }\}$$

where   *A(q) = { a 2 A : (9 q') (q,a,q') 2 T}*

Note: I/O automata are input enabling ($A^I(q) = A^I$ for all q) and therefore have no error states.

ERROR state of the product.

The Compatibility of Interface Automata

An *environment* for an interface automaton *F* is an interface automaton *E* such that

     1. *F* is composable with *E*

     2. $A^I_F = A^O_E$

     3. *Error(F,E) = ;*

A *helpful* environment for two composable interface automata *F* and *G* is an environment *E* for the product *F £ G* such that

     4. no state in *Error(F,G) £ Q_E* is reachable in *(F £ G) £ E*

Two interface automata *F* and *G* are *compatible* if they are composable and there exists a helpful environment.

ok

ack?

ack?

msg send! nack? send!

nack?

send

ack nack

send ack nack

send?

ack!

Environment can avoid this state.

The most general helpful environment.

Computing the Composite Interface Automaton

1. Construct product automaton.

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1. Construct product automaton.

2. Mark error states as incompatible.

Computing the Composite Interface Automaton

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3. Until no more incompatible states can be added: mark state q as incompatible if the environment cannot prevent an incompatible state to be entered from q.

Computing the Composite Interface Automaton

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2. Mark error states as incompatible.

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4. If the initial state is incompatible, then the two interfaces are incompatible. Otherwise, the composite interface is the product automaton without the incompatible states.
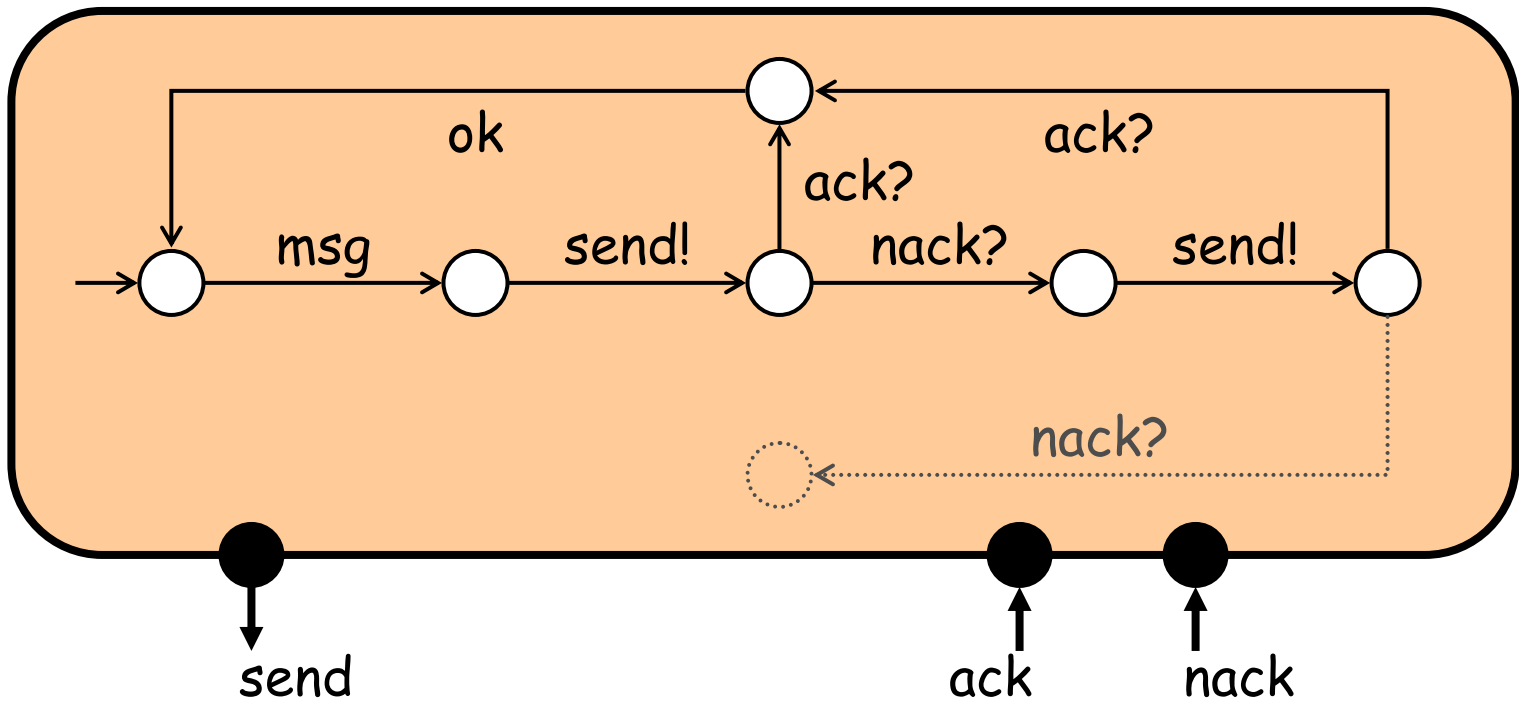
Computing the Composite Interface Automaton

1. Construct product automaton.

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3. Until no more incompatible states can be added: mark state q as incompatible if the environment cannot prevent an incompatible state to be entered from q.

4. If the initial state is incompatible, then the two interfaces are incompatible. Otherwise, the composite interface is the product automaton without the incompatible states.

This procedure computes the most general helpful environment as the most general strategy of the environment to avoid error states.

The composite interface automaton.